

A New Framework for Scalable Secure Block Cipher Generation using Color Substitution and Permutation on Characters, Numbers, Images and Diagrams

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ABSTRACT

The aim of our research is to develop a new framework for secure block cipher generation using color substitution and permutations on alphanumeric letters, symbols, images, diagrams or any kind of text. To transfer the keys from source to destination we have used RSA public key algorithm and for encryption / decryption of the information, we have used our invented ‘play color substitution’ algorithm. Importance of RTF will be explained, Cryptanalysis attacks were discussed and shown that the cipher cannot be broken by any cryptanalysis attacks.

General Terms

Crypt analysis, block cipher, play color cipher, decillions, encryption / decryption algorithm.

Keywords

RSA, ETF: Electronic Frontier Foundation, PUB: Public key of user B, PRA: Private Key of user A, PUA: Public key of user A, PRB: Private key of user B, PCC: Play color cipher, RTF: Rich text format, LHS: Left hand side, RHS: Right hand side.

1. INTRODUCTION

The most prevailing and universal approach to countering the threats to network / data security is encryption. Even though it is very influential, the cryptanalysts are very intellectual and they were working round the clock to break the ciphers.

Many Scholars and Scientist were investigating day and night on the existing algorithms to make more stronger and unbreakable ciphers by enhancing them[1,2,3,4,5,6,7,8,9,10]. But still most of the algorithms were vulnerable to attack. One of the most widely used cryptographic algorithm is DES, is also broken and announced by the Electronic Frontier Foundation in July 1998 [16], many substitution algorithms like Caesar Cipher, Mono alphabetic Cipher, Play fair Cipher and Poly alphabetic Ciphers are proven to be not stronger and even they only support for limited applications.

To meet the current requirements and to fight with the cryptanalyst in the battle of the network and information security, we need stronger Cryptographic algorithms.

It is mandatory that to sharp our brain and invent new Cryptographic algorithms rather than updating the existing one all the time. In lieu of this and to bring the revolution in the field of network and Information security; have invented

a new cryptographic algorithm by using ‘Color substitution technique’ and named it as ‘Play color Cipher’ in our previous paper [15].

In the present paper we have used 92 bit key and extended the algorithm to the wide range of applications, where it allows to encrypt / decrypt all types of characters, numbers, symbols, diagrams and images, etc.. To create stronger cipher we did permutations/ transposition on the output of R.T.F and before applying color substitution. Finally, we have proven that the cipher is very stronger and resistant to the cryptanalysis attacks.

2. KEY SELECTION & DISTRIBUTION

Sequence of steps involved in *selecting a key and procedure for transferring key from source to the destination* as shown in the Figure 1 and Figure 2:

- Select key ‘K’, should be 23 decimal numbers between 0 to 9 (having 3 sub keys)
The first 14 digits in the Key can be between 0000000000 0001 (Min) to 9999999999 9999 (Max)
Remaining 9 digits (RHS) of the key should be the numbers between 1 to 9, and the number once used should not be repeated.
- In the above 23 decimal numbers:, from LHS to RHS, algorithm considers first 10 numbers as starting address (K1), next 4 numbers as increment value(K2) and the last 9 numbers as key (K3)for transposition / permutation..
- Use RSA [16] Public key encryption algorithm for key distribution as shown in Figure 2:
- Encrypt K using receivers (User B) Public key (PUB) for confidentiality ----- 2.1
- Encrypt the result of 2.1 using senders (User A) Private key (PRA) for Authentication. ----- 2.2
- Send the result of 2.2 to the receiver-----2.3
- Decrypt 2.3 by using PUA ----- 2.4
- Decrypt 2.4 by using PRB ----- 2.5

Hence with both authentication and confidentiality we have distributed the keys between User A and User B.

3. METHODOLOGY

Sequence of steps in our ‘Play color cipher’ algorithm for encryption and decryption:

Procedure for encrypting the plaintext at the sender as shown in Figure 3,4,6 and 7:

- Use Windows XP and Microsoft Visual Studio for executing “*Play Color Cipher Algorithm*”.
- Debug the algorithm to get the ‘*Form for Encryption*’

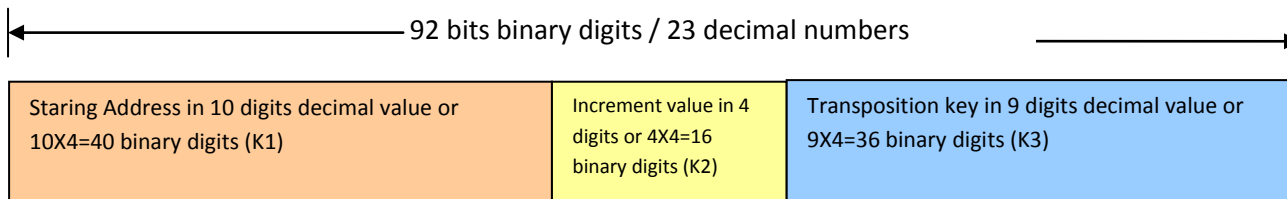


Figure 1: Key format in 92 binary bits

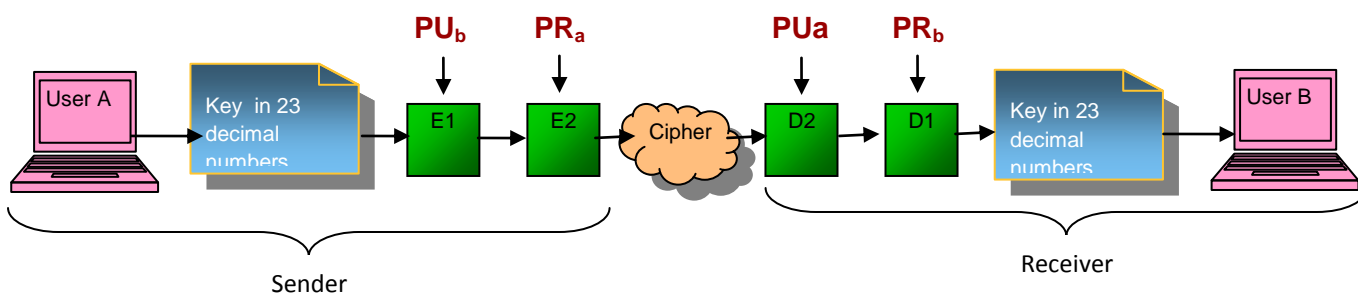


Figure 2: Secure Transmission of Key using RSA

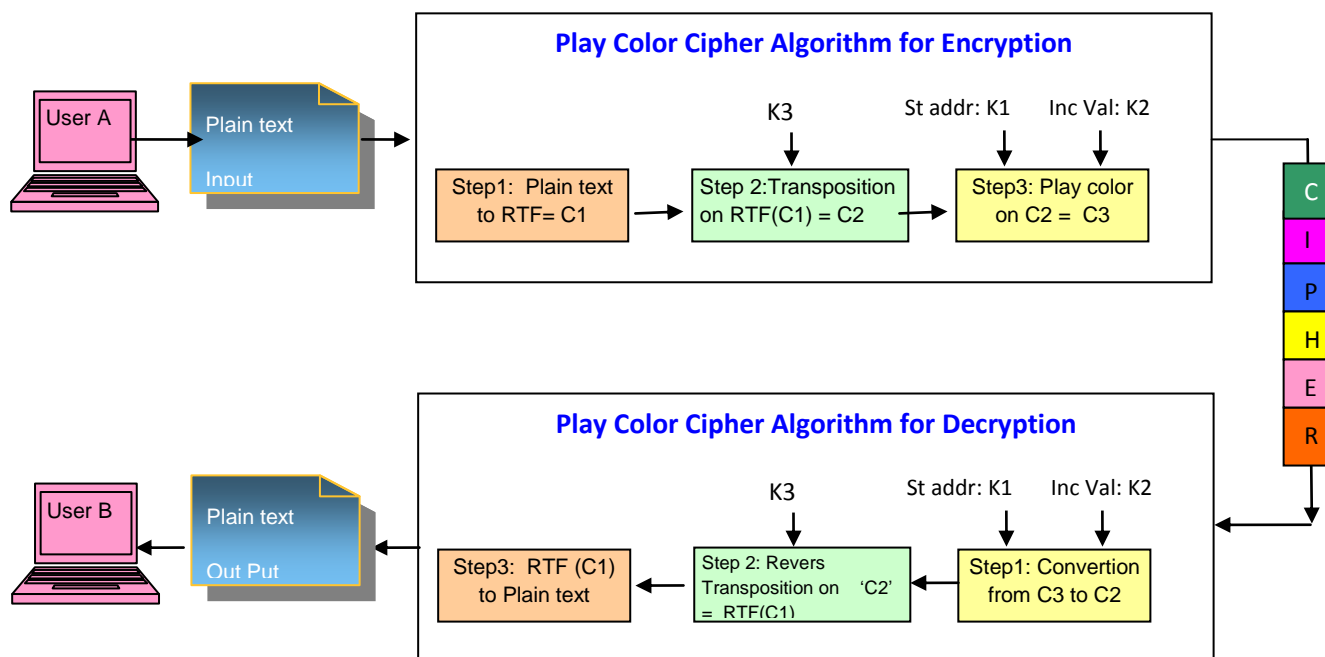


Figure 3: Process of Encryption & Decryption using Play Color Cipher

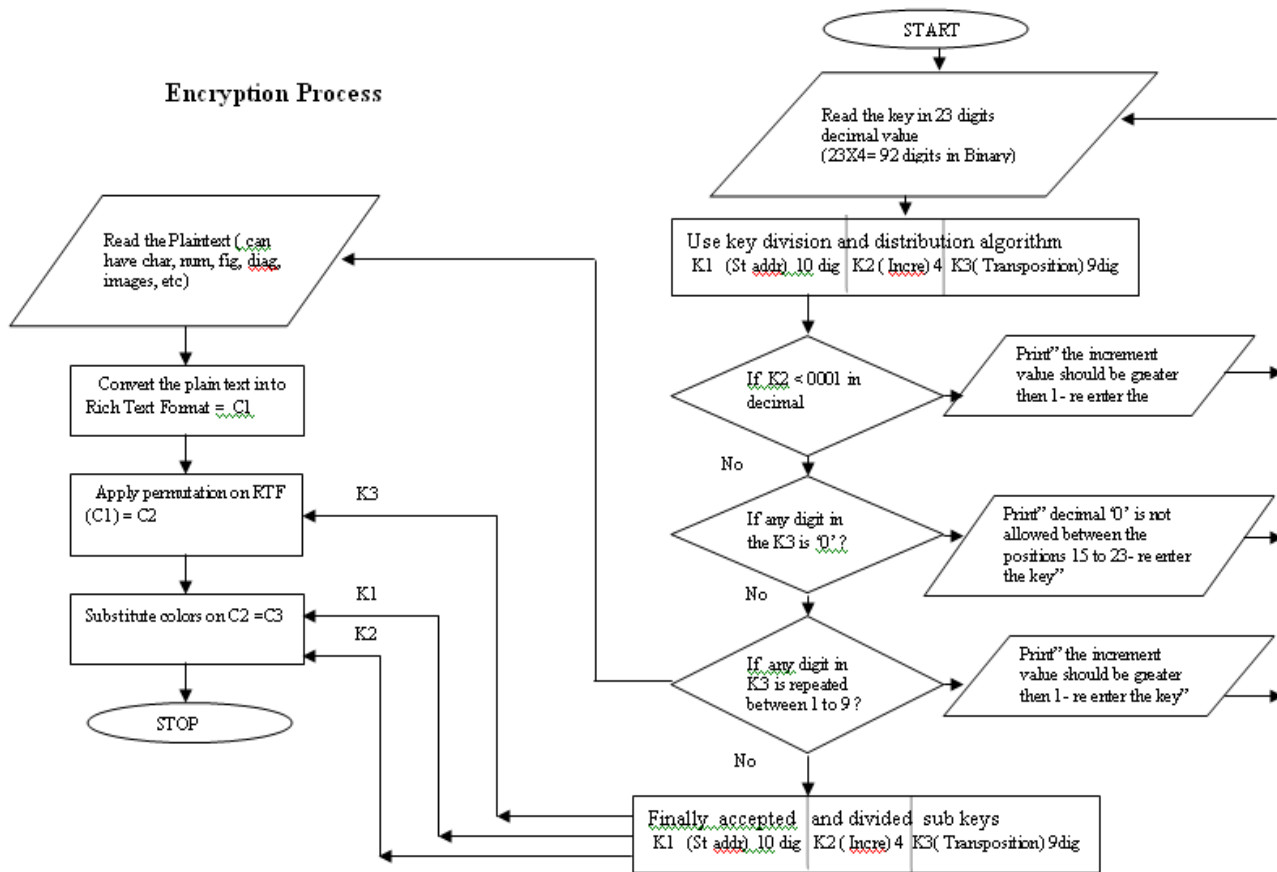


Figure 4: Flow chart for encryption process

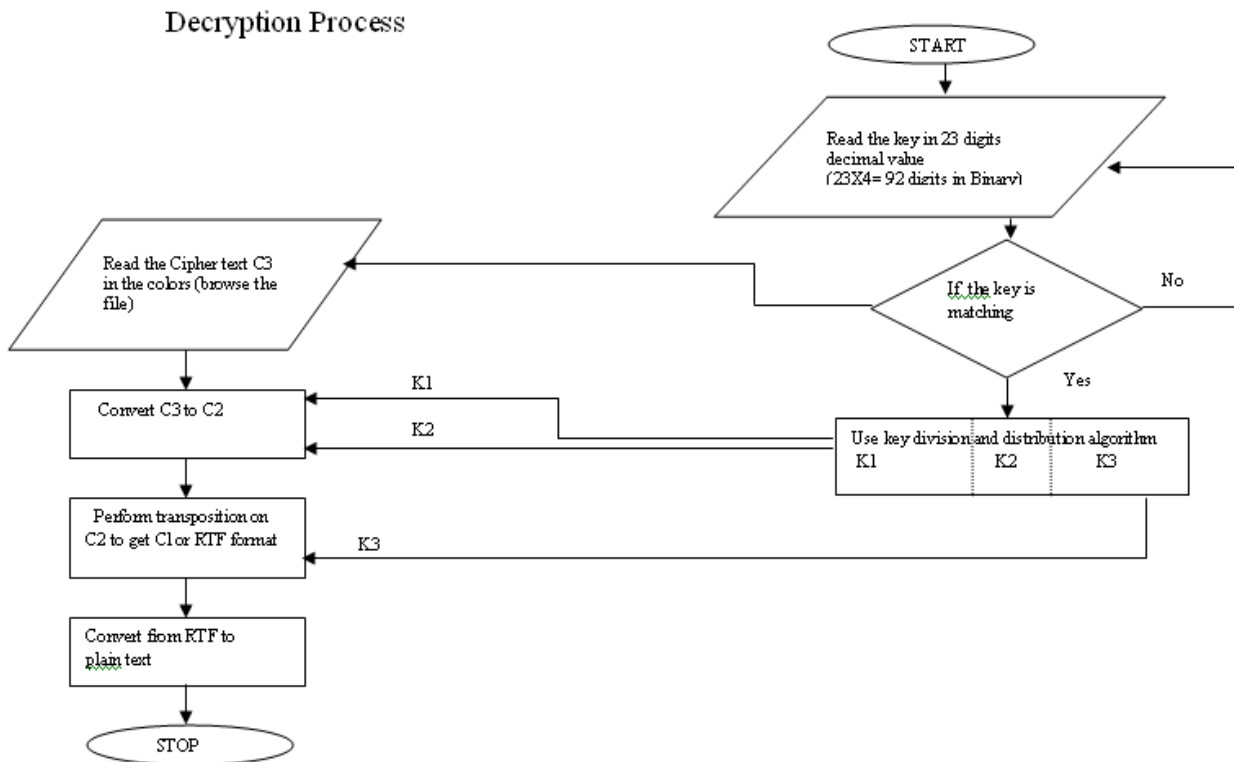


Figure 5: Flow chart for Decryption Process

- Enter the values of key ‘K’ in the box given and click on ‘assign button’
- Enter the plain text in the box given (can use any characters, numbers, symbols, figures, diagrams or any kind of text).
- Click on ‘Encryption Button’
- ‘Save’ the result as a file with the extinction of .tif and ‘exit’ from encryption form.
- Now the cipher is ready for transferring from source to destination.

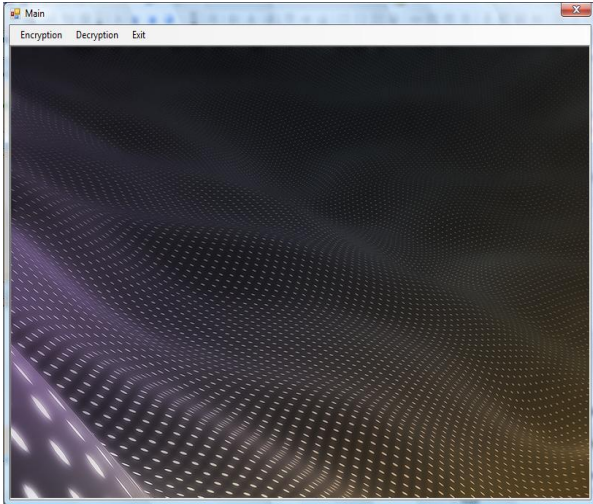


Figure 6: Main form of ‘play color cipher algorithm’

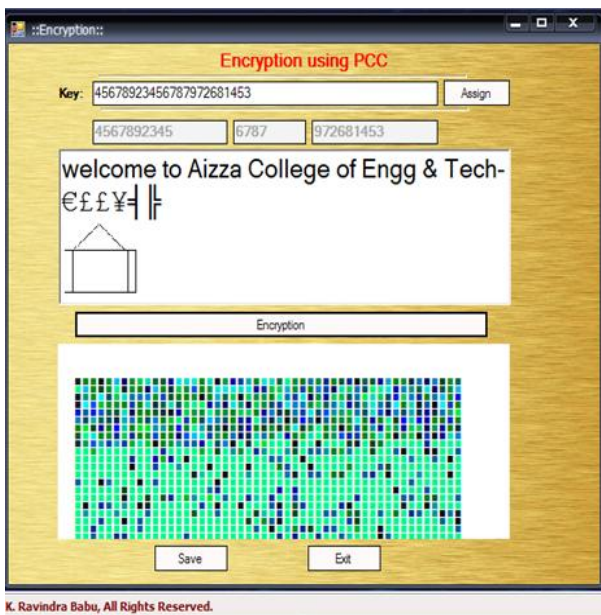


Figure 7: Form for Encryption

Procedure for Decrypting the cipher text at the receiving end as shown in Figure 3,5 and 8:

- Use Windows XP and Microsoft Visual Studio for executing ‘Play Color Cipher Algorithm’.
- Debug the algorithm to get the ‘Form for Decryption’
- Open the ‘Form for Decryption’ and enter the ‘K’
- Browse for the received and saved cipher file
- Click on ‘Decryption Button’
- Observe the converted ‘plain text’.

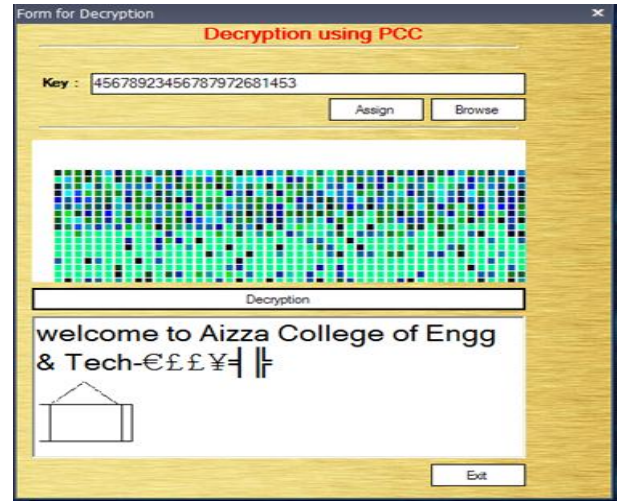


Figure 8: Form for Decryption


4. DEVELOPMENT OF THE CIPHER

4.1 About RTF and Converting a Plain text in to the RTF format:

About Rich Text Format (RTF): The Windows Forms Rich Text Box control is used for displaying, entering, and manipulating text with formatting. The Rich Textbox control does everything the Text Box control does, but it can also display fonts, colors, and links; load text and embedded images from a file; and find specified characters. The Rich Textbox control is typically used to provide text manipulation and display features similar to word processing applications such as Microsoft Word. We can convert all types of characters, numbers, symbols and diagrams by using rich text box in to Rich text format. By using this we can convert the plaintext into an unintelligible text. The output of this step we name it as Cipher text C1.

Consider the plain text as shown below, it may contain characters, numbers, symbols, diagrams, images e.t.c., by using rich text box it is converted in to an un intelligible form as shown below; it is noticeable that the diagrams or the images in the plain text is also got converted into numbers and symbols.

Plain text:

welcome to Aizza College of Engg & Tech-
€££¥|=|


Rich text / Cipher text ‘C1’ for the above plain text:

```

\rtf1\fbidis\ansi\ansicpg1252\deff0\deflang1033{\fonttbl{\f0\fnil\fcharset0 Microsoft Sans
Serif;}{\f1\modern\prq1\fcharset0 Courier New;}}
\viewkind4\uc1\pard\ltrpar\fo\fs20 welcome to Aizza College of Engg & Tech
\lang1024\fi\fs24\80\as\as5\us95692\us9568? \lang1033\fo\fs20\par
\pic1\wmetafile8\picw1931\pich2752\picwgoal1095\pichgoal1560
010009000003cc05000006004a04000000004a04000026060f08a08574d46430100000000000001
007cb9000000000100000680800000000005808000001000006c000000ffffff000000000
4a000000680000000000000000000004b060000fe080000204544600000100680800004c000f
0003000000000000000000000000000004000000030000e2000000aa0000000000000001
000000000000d72030010980200160000000c00000018000000a0000001000000000000000
b00000090000001000000490000006800000052000007001000001000000f1ffffff00000000
00000000000009001000000000004400022430061006c006900620072006900000000000000
00000000000000000000000000000000000000000000000000000000000000000000000000
0000000000000001300d8d91300100000003cdd1300bcda130035060323cdd130034da13001f
    
```


4.2 Performing Permutation on the output of Previous step (4.1) using K3

To make stronger cipher from the previous step 4.1, we have performed permutation on C1 by using key K3. The sub key K3 is a 9 digits decimal number as shown in the Figure 1,4. The numbers in the K3 can be between 1 to 9, zero is not allowed to use and the number once used should not be repeated.

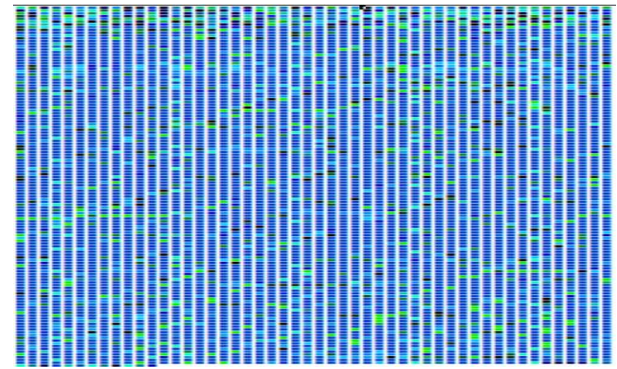
For performing transposition write the message in the rectangle, row by row, and read the message off, column by column, but permute the order of the columns. The order of the columns then becomes key to the algorithm. In the example shown below, the key- K3 is '972681453' used for performing transposition on the cipher C1, the output is as follows and named it as Cipher C2:

```
b(\rtf1\fiidisans1\ansicpg0252\deff1\deflangt033\fonttbl{\f0\rl\vfcharset0 Microsoft Sa);
Serif.d\fi\fmolern\forat\vfcharser0 Courier
New;)}
d\viewkinnr4\uc1\pa\d\trparwfo\fs20 oelSome to Aizee C llaag ofeEngg & T1ch
\lang024\fa24\80\53\as3\au\95692\9568? \0ng1033\fa\fs20\pa\
\pictewmetafil38\picw1951\picw27a2\picwgc1095\pi0hgoal1569
010000000003cc05000006004a04000000004a04000026060d08a085740464301001000000000
007cb91000000000000000680800000000000068080000010000006c0000f0fffff
000000000
4a000000680000000000000000000000004b860000fe04000020451d4600000000680800
04c0000
0000300000000000000000000000000000004000000030000e200000aa0000000000000000
000000000000d072030010980200160000008c00000010000000a00000010000000000000000
000000090000000100000004900000068000007520000010010000f000000f0fffff00000000
000000000000000090010000000020004400002430061066c0069009200720060000000000000
00000000000000000000000000000000000000000000000000000000000000000000000000
000000000000d001300d8091300100100003cdd3300bcda1300035060023cdd130134da130000
000002a4db130090dd13006d4e60323cad130034d013002000f0004964203134da13003cdd1300
200000f0fffff0c02e000f0642f31ffffffffff0180ff00180edffff180fffff
ff0000000
```

4.3 Applying play color cipher on the output of previous step(4.2) using K1and K2

Play Color Cipher: Each Character (Capital, Small letters, any kind of text, Numbers (0-9), Symbols) in the plain text is substituted with a color block from a 18 decillions of colors[15] available in the computer world. In this we have considered only ARGB with the maximum number of 255 X 255 X 255 X 255 = 4228250625 colors, to make the cipher more stronger we have used starting address of the color K1 in 10 digits decimal number from 0000000000 to 9999999999 and increment value K2 in 4 digits decimal number with the range 0001 to 9999. It is to be noted that increment value should not be 0000 as shown in the Figure 4, because all the characters will get the same color, which is not acceptable in algorithm.

In the example shown below the value of the starting address K1 is '4567892345' and the Increment value K2 is ' 6787', by applying this on the output Cipher(C2) in previous step, we got the color code as shown below. This is the final cipher C3 generated by the source/ sender intended for destination/ receiver.



5. CRYPTANALYSIS

The cryptanalyst attacks which are generally considered in the literature of Cryptography are

1. Cipher text only attack (Brute force attack)
2. Known plaintext attack
3. Chosen plaintext attack
4. Chosen cipher text attack

In this analysis the key K is consisting of 23 decimal numbers, where in each number can be represented in the form of 4 binary bits. Hence the length of the key is 92 bits and the size of the key space is

$$2^{92} = 5.0 \times 10^{27} \text{ Keys}$$

If the time required for the determination of the plain text for one value of the key in the key space is taken as 10^{-3} seconds, then the time required for obtaining the plain text by considering all the possible keys in the key space is

$$\frac{5.0 \times 10^{27} \times 10^{-3}}{365 \times 24 \times 60 \times 60} = 1585 \times 10^{14} \text{ Years}$$

If we perform one encryption per micro second it takes 2.4×10^{16} years, and for 10^6 encryptions per micro second it leads to 1842.6 years.

This number is very large; hence, it is impracticable to break the cipher.

In the case of known plain text attack, we know as many pairs of plaintext and cipher text as we require. The number of colors in the computer world is more than 18 decillions, with minor difference we have thousands of shades in the same color, by looking at the colors it is impossible to obtain the plain text, even if you have number of plain text and the corresponding cipher text, the plain text is not the exact plain text of the color cipher because in step one we have converted the plain text in to RTF format, considered the result as C1, then in the second step we have performed trans position on the C1 and the result in this step is C2, in third step we did color substitution on the C2 and the result is C3. Hence, the plain text for final cipher C3 is another cipher C2, but not the exact plain text. With this permutations and substitutions in different stages we can conclude that knowing plain text does not work.

In the last two cases of the cryptanalysis attack, no scope is found for breaking the cipher.

In view of the above discussion, we conclude that the Cipher is a very strong.

6. RESULTS

The invented play color cipher algorithm works with 92 bit key and it is proven that it is comfortably converting all kinds of text, symbols, diagrams and images. The process of conversion with example was explained in previous section. The strength of the any algorithm depends on key rather than the algorithm, in this the length of the key is 23 decimal digits and proven that it is far from crypt analysis attacks.

7. CONCLUSION

In this paper we have presented a conventional encryption scheme using color substitution and permutations with symmetric key. The length of the block is variable in size. we have proven that it can encrypt / decrypt all kinds of text, numbers, symbols, images and diagrams with example. Algorithm have three functions, works with three sub keys K1, K2 and K3 generated from the main key K. Using the existing 'Rich text box' feature from Microsoft Visual studio is the center of attraction of this algorithm. For transferring key from source to destination we have used RSA algorithm and the procedure was explained with neat diagram.

Applying the combination of permutation and substitution enhanced the strength of the algorithm to the great extent. This algorithm is implemented by using C#.net. The brief explanation and the advantages of RTF were given; Generation of cipher text in three stages was explained with example. With the 92 bit key the cipher is very strong and far from cryptanalyst attacks. For performing 10^6 encryptions per micro second it takes 1842.6 years. Finally we conclude that the algorithm is potential one.

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